# A SHORTEST PATH LENGTH ON A FUZZY NETWORK WITH TRIANGULAR INTUITIONISTIC FUZZY NUMBER 

A. D. Chandrasekaran, S. Balamuralitharan and K. Ganesan<br>Faculty of Engineering and Technology, Department of Mathematics, SRM University, Kattankulathur, Tamilnadu, India<br>E-Mail: adchandru1977@gmail.com


#### Abstract

In this paper, a fuzzy network method is proposed for finding the Shortest Path length (SPL) with Triangular Intuitionistic Fuzzy Numbers (TIFN). Furthermore, to find the smallest path of the edge by the intuitionistic fuzzy distance using graded mean integration. We discussed the SPL from a specified vertex to all other edges in a fuzzy network. An illustrative example is given to express our proposed work.


Keywords: triangular intuitionistic fuzzy number, fuzzy network, shortest path length, bellman's equation.
AMS classifications: 03B52, 03E72, 05C72, 90C70.

## 1. INTRODUCTION

In this paper, we proposed a new network approach that can obtain the triangular intuitionistic fuzzy number. The SPL requires determining the shortest route for fuzzy network. The SPL is one of the most fundamental and well-known combinatorial optimization problems that appear in many applications as a subproblem.

In recently, many researchers have developed much concentration to the fuzzy SPL since it is essential to many of applications [1-3, 6-10]. In some applications, the numbers associated with the edges of networks may represent characteristics other than lengths, and we may want the optimum paths, where optimum can be defined by different criteria. In the fuzzy SPL, the fuzzy network shortest length and the corresponding shortest path problems are the useful information for the decision makers. The length of arcs in the network represents travelling time, cost, distance or other variables.

The SPL in a fuzzy network has discussed many previous researchers since it is important to many applications such as communications, routing and transportation $[4,12,13]$. In a fuzzy network problem, the arcs length is assumed to represent transportation cost. In the real life, the transportation cost may be known only approximately due to pensiveness of information.

The fuzzy SPL was first developed by Dubois and Prade [5]. Although the shortest path length can be obtained, a corresponding shortest path cannot be identified [11]. Liu and Kao's algorithm for finding the shortest path can obtain a non-dominated shortest path [14]. In1980 Dubois and Prade [5] first introduce fuzzy SPP. Okada and Soper [11] developed an algorithm based on multiple labeling approaches which is useful to generate number of non-dominated paths. Applying minimum concept they have introduced an order relation between fuzzy numbers. Applying extension principle Klein [4], has given an algorithm which results in a dominated path on a fuzzy network.

In this paper is organized as follows. In section 2, preliminary ideas and definitions are given. The procedure for finding SPL using TIFN derived in section 3 and 5. An
example is provided in section 4 to find the shortest path length and discussed the conclusion in section 5.

## 2. PRELIMINARIES

## Definition 2.1

Let $T=\left\{x, \mu_{T}(x), \gamma_{T}(x) / x \in X\right\}$ be an IFS, then we call $\left(\mu_{T}(x), \gamma_{T}(x)\right)$ an Intuitionistic Fuzzy Number (IFN). We denote it by $(\langle x, y, z\rangle,\langle l, m, n\rangle)$ where $\langle x, y, z\rangle$ and $\langle l, m, n\rangle \in F(I), I=[0,1], 0 \leq z+n \leq 1$.

## Definition 2.2

An Intutitionistic Fuzzy Set (IFS) T in X is given by $T=\left\{x, \mu_{T}(x), \gamma_{T}(x) / x \in X\right\}$ where $\mu_{T}(x): X \rightarrow[0,1]$ and $\quad \gamma_{T}(x): X \rightarrow[0,1]$ and for every $x \in X, 0 \leq \mu_{T}(x)+\gamma_{T}(x) \leq 1$.

## Definition 2.3

|  | A Triangular Intuitionistic Fuzzy |  |
| :---: | :---: | :---: |
| ' | is |  |
| $A=(\langle x, y, z\rangle,\langle l, m, n\rangle)$ with $\langle l, m, n\rangle \leq\langle x, y, z\rangle^{c}$ that |  |  |

either $1 \geq y, m \geq z$ (or) $m \leq x, n \leq y$ are membership and non-membership fuzzy numbers of A .

The additions of two TIFN are as follows:
For two triangular intuitionistic fuzzy numbers
$T_{1}=\left(\left\langle x_{1}, y_{1}, z_{1}\right\rangle: \mu_{T_{1}},\left\langle l_{1}, m_{1}, n_{1}\right\rangle: \gamma_{T_{1}}\right)$ and
$T_{2}=\left(\left\langle x_{2}, y_{2}, z_{2}\right\rangle: \mu_{T_{2}},\left\langle l_{2}, m_{2}, n_{2}\right\rangle: \gamma_{T_{2}}\right)$
with $\quad \mu_{T_{1}} \neq \mu_{T_{2}}$ and $\quad \gamma_{T_{1}} \neq \gamma_{T_{2}}$, define
$T_{1}+T_{2}=\binom{\left\langle x_{1}+x_{2}, y_{1}+y_{2}, z_{1}+z_{2}\right\rangle: \operatorname{Min}\left(\mu_{T_{1}}, \mu_{T_{2}}\right)}{,\left\langle l_{1}+l_{2}, m_{1}+m_{2}, n_{1}+n_{2}\right\rangle: \operatorname{Max}\left(\mu_{T_{1}}, \mu_{T_{2}}\right)}$.
www.arpnjournals.com

## Definition 2.4

For each $S_{d}=(\langle a, b, c\rangle,\langle e, f, g\rangle) \in F$, the signed distance of $S_{d}$ measured from 0 and is defined by $d\left(S_{d}, 0\right)=\frac{1}{4}[(a+e), 2(b+f),(c+g)]$. from this definition for each $S_{d}=\left[\left(\left\langle a_{1}, a_{2}, a_{3}\right\rangle,\left\langle a_{4}, a_{5}, a_{6}\right\rangle\right),,\left\langle\left\langle b_{1}, b_{2}, b_{3}\right\rangle,\left\langle b_{4}, b_{5}, b_{6}\right\rangle\right),,\left\langle\left\langle c_{1}, c_{2}, c_{3}\right\rangle,\left\langle c_{4}, c_{5}, c_{6}\right\rangle\right)\right]$, we obtain $s_{d}=\frac{1}{4}[A+2 B+C]$, where
$A=\left(a_{1}+a_{4}, a_{2}+a_{5}, a_{3}+a_{6}\right), B=\left(b_{1}+b_{4}, b_{2}+b_{5}, b_{3}+b_{6}\right)$ and $C=\left(c_{1}+c_{4}, c_{2}+c_{5}, c_{3}+c_{6}\right)$.

## 3. SHORTEST PATH LENGTH (SPL) PROCEDURE IN BELLMAN'S EQUATION

According to Bellman's equation, a Dynamic Programming formulation for the SPL can be given as follows:

Given a network with an acyclic directed graph $G=(V, E)$ with n vertices numbered from 1 to n such that 1 is the source and n is the destination. Then we have,

$$
g(n)=0, g(i)=\min _{i<j}\left\{e_{i j}+g(j) /\langle i, j\rangle \in E\right\} .
$$

Here $e_{i j}$ is the weight of the directed edge $\langle i, j\rangle$ and $g(i)$ is the length of the shortest path from vertex $i$ to vertex $n$. From the following figure, the solution of Dynamic Programming can be derived as follows:

## 4. AN EXAMPLE



Consider a network with the triangular intuitionistic fuzzy arc lengths as shown below. The arc lengths are assumed to be
$E_{12}=(\langle 17,30,42\rangle,\langle 37,50,56\rangle) ;$
$E_{13}=(\langle 25,35,49\rangle,\langle 37,51,56\rangle)$;
$E_{23}=(\langle 33,49,56\rangle,\langle 52,61,69\rangle)$;
$E_{24}=(\langle 33,41,53\rangle,\langle 42,55,64\rangle)$;
$E_{35}=(\langle 14,21,29\rangle,\langle 24,31,39\rangle)$;
$E_{45}=(\langle 14,38,52\rangle,\langle 44,52,54\rangle)$;
$E_{46}=(\langle 20,28,36\rangle,\langle 30,40,49\rangle) ;$
$E_{56}=(\langle 23,35,52\rangle,\langle 43,55,65\rangle)$.
The possible paths are as follows:
$W_{1}: 1-2-4-5-6 ; \mathrm{L}_{1}=(\langle 87,144,199\rangle,\langle 166,212,239\rangle)$
$W_{2}: 1-2-4-6 ; \mathrm{L}_{2}=(\langle 70,99,131\rangle,\langle 109,145,169\rangle)$
$W_{3}: 1-2-3-5-6 ; \mathrm{L}_{3}=(\langle 87,135,179\rangle,\langle 156,197,229\rangle)$
$W_{4}: 1-3-5-6 ; \mathrm{L}_{4}=(\langle 62,91,130\rangle,\langle 104,137,160\rangle)$.
The solution of Dynamic Programming can be derived as follows:

$$
\begin{aligned}
g(6) & =0, g(5)=E_{56}=(\langle 23,35,52\rangle,\langle 43,55,65\rangle) \\
g(4) & =\min _{4<j}\left\{e_{4 j}+g(j) /\langle 4, j\rangle \in E\right\} \\
& =\min \left\{E_{46}, E_{45}+E_{56}\right\} \\
& =E_{46}=(\langle 20,28,36\rangle,\langle 30,40,49\rangle) \\
g(3) & =\min _{3<j}\left\{e_{3 j}+g(j) /\langle 3, j\rangle \in E\right\} \\
& =E_{35}+g(5) \\
& =E_{35}+E_{56}=(\langle 37,56,81\rangle,\langle 67,86,104\rangle)
\end{aligned}
$$

Similarly, we obtain,

$$
\begin{aligned}
g(2) & =\min \left\{E_{24}+g(4), E_{24}+E_{25}+g(5)\right\} \\
& =E_{24}+E_{46}=(\langle 53,69,89\rangle,\langle 72,95,113\rangle)
\end{aligned}
$$

and
$g(1)=\min \left\{E_{12}+g(2), E_{13}+g(3), E_{12}+E_{23}+g(3), E_{12}+E_{24}+g(4), E_{12}+E_{24}+E_{45}+g(5)\right\}$ $=(\langle 62,91,130\rangle,\langle 104,137,160\rangle)$.

## 5. TO COMPUTE THE SHORTEST PATH LENGTH (SPL)

In this problem we consider is that the edge weight in the network denoted by $e_{i j}$ and the edge weight should be expressed using fuzzy linguistics, and also this used in TIFN.
$\tilde{e}_{i j}=\left(\alpha, e_{i j}, \beta\right) \cdots \cdots(2)$
where $\alpha=e_{i j}-\delta_{i j}, \beta=e_{i j}+\delta_{i j}$ and $0<\delta_{i j}<e_{i j}, \delta_{i j}>0$ since $\delta_{i j}$ should be determined by the Decision Maker. Now consider the fuzzy case. We look for inequalities that satisfy $\quad e_{i i_{1}}+e_{i_{1} i_{2}}+\cdots+e_{i_{m}(i)^{n}} \leq e_{i j_{1}}+e_{j_{1} j_{2}}+\cdots+e_{j_{m}(i)^{n}}$, $\forall i<j,\langle i, j\rangle \in E$. when $i=1, e_{13}+g(3)<e_{12}+g(2)$, that is
$e_{13}+e_{35}+e_{56}<e_{12}+e_{23}+g(3)<e_{12}+e_{23}+e_{35}+e_{56}$ (or)
$e_{13}+g(3)<e_{12}+e_{24}+e_{45}+e_{56}$.
when $i=2, e_{24}+g(4)<e_{23}+g(3)$, that is $e_{24}+e_{46}<e_{23}+e_{35}+e_{56}$.
when $i=4, e_{46}+g(6)<e_{45}+g(5)$, that is $e_{46}<e_{45}+e_{56}$.
www.arpnjournals.com

Then
$\delta_{i i_{1}}+\delta_{i_{1} i_{2}}+\cdots+\delta_{i_{m}(i)^{n}} \leq \delta_{i j_{1}}+\delta_{j_{1} j_{2}}+\cdots+\delta_{j_{m}(i)^{n}}$,
$\forall i<j,\langle i, j\rangle \in E$ based on the above inequalities are derived as

$$
\begin{align*}
\delta_{13}+\delta_{35}+\delta_{56} & <\delta_{12}+\delta_{24}+\delta_{46}<\delta_{12}+\delta_{24}+\delta_{45}+\delta_{56} \cdots \cdots(3)  \tag{3}\\
\delta_{24}+\delta_{46} & <\delta_{23}+\delta_{35}+\delta_{56} \\
\delta_{46} & <\delta_{45}+\delta_{56}
\end{align*}
$$

| $\tilde{e}_{i j}$ | $\alpha=e_{i j}-\delta_{i j}$ | $e_{i j}$ | $\beta=e_{i j}+\delta_{i j}$ |
| :---: | :---: | :---: | :---: |
| $\tilde{e}_{12}$ | $(\langle 14,26,36\rangle,\langle 32,43,48\rangle)$ | $(\langle 17,30,42\rangle,\langle 37,50,56\rangle)$ | $(\langle 20,34,48\rangle,\langle 42,57,64\rangle)$ |
| $\tilde{e}_{13}$ | $(\langle 20,29,42\rangle,\langle 29,42,46\rangle)$ | $(\langle 25,35,49\rangle,\langle 37,51,56\rangle)$ | $(\langle 30,41,56\rangle,\langle 45,60,66\rangle)$ |
| $\tilde{e}_{23}$ | $(\langle 29,41,46\rangle,\langle 43,50,57\rangle)$ | $(\langle 33,49,56\rangle,\langle 52,61,69\rangle)$ | $(\langle 37,57,66\rangle,\langle 61,72,81\rangle)$ |
| $\tilde{e}_{24}$ | $(\langle 30,35,46\rangle,\langle 35,46,54\rangle)$ | $(\langle 33,41,53\rangle,\langle 42,55,64\rangle)$ | $(\langle 36,47,60\rangle,\langle 49,64,74\rangle)$ |
| $\tilde{e}_{35}$ | $(\langle 12,17,23\rangle,\langle 19,24,31\rangle)$ | $(\langle 14,21,29\rangle,\langle 24,31,39\rangle)$ | $(\langle 16,25,35\rangle,\langle 29,38,47\rangle)$ |
| $\tilde{e}_{45}$ | $(\langle 10,32,45\rangle,\langle 37,43,43\rangle)$ | $(\langle 14,38,52\rangle,\langle 44,52,54\rangle)$ | $(\langle 18,44,59\rangle,\langle 51,61,65\rangle)$ |
| $\tilde{e}_{46}$ | $(\langle 14,20,27\rangle,\langle 20,29,36\rangle)$ | $(\langle 20,28,36\rangle,\langle 30,40,49\rangle)$ | $(\langle 26,36,45\rangle,\langle 40,51,62\rangle)$ |
| $\tilde{e}_{56}$ | $(\langle 19,29,45\rangle,\langle 34,44,53\rangle)$ | $(\langle 23,35,52\rangle,\langle 43,55,65\rangle)$ | $(\langle 27,41,59\rangle,\langle 52,66,77\rangle)$ |

From (2.4), we obtain the following estimate of the edge weights in the fuzzy sense:

| $E_{12}^{0}=\langle 54,80,98\rangle$ | $E_{13}^{0}=\langle 62,86,105\rangle$ | $E_{23}^{0}=\langle 85,110,125\rangle$ |
| :---: | :---: | :---: |
| $E_{24}^{0}=\langle 75,96,117\rangle$ | $E_{35}^{0}=\langle 38,52,68\rangle$ | $E_{45}^{0}=\langle 58,90,106\rangle$ |
| $E_{46}^{0}=\langle 50,68,85\rangle$ | $E_{56}^{0}=\langle 66,90,117\rangle$ |  |

The fuzzy network $G=(V, E)$ with $\left\{e_{i j}^{0} /\langle i, j\rangle \in E\right\}$.

## where

$g^{0}(1)=E_{13}^{0}+g(3)=E_{13}^{0}+E_{35}^{0}+g(5)=E_{13}^{0}+E_{35}^{0}+E_{56}^{0}$.
The fuzzy shortest path is $1-3-5-6$ with length $\langle 166,228,290\rangle$.

## 6. CONCLUSIONS

In this paper, a shortest path length is obtained using a procedure based on Bellman's equations in a fuzzy network. The SPL arc lengths are considered as uncertain and are characterized by triangular intuitionistic fuzzy numbers. The fuzzy shortest length is the useful information for the decision makers in an intuitionistic fuzzy SPL. This algorithm can be computed using the TIFN by the decision maker. It provides the good output for different types of fuzzy network. The effectiveness of the method is tested by an example.

## REFERENCES

[1] A. Kiran Yadav, B. Ranjit Biswas. 2009. Finding a Shortest Path using an Intelligent Technique. International Journal of Engineering and Technology. 1(2): 1793-8326.
[2] A. Nagoor Gani and M.Mohamed Jabarulla. 2010. On Searching Intuitiionistic Fuzzy Shortest Path in a Network. Applied Mathematical Sciences. (29): 34473454.
[3] Amit Kumar and Manjot Kaur. 2011. A new algorithm for solving network flow problems with fuzzy arc lengths. An official journal of Turkish Fuzzy systems Association. 2(1): 1-13.
[4] C.M.Klein. 1991. Fuzzy Shortest paths. Fuzzy Sets and Systems. 39(1): 27-41.
www.arpnjournals.com
[5] D. Dubois and H.Prade. 1980. Fuzzy Sets and Systems: Theory and Applications, Academic Press, New York.
[6] J.S. Yao and K.M. Wu. 2000. Ranking Fuzzy numbers based on decomposition principle and signed distance. Fuzzy Sets and Systems. 116: 275-288.
[7] K.C.Lin and M.S. Chern. 1993. The Fuzzy shortest path problem and its most vital arcs. Fuzzy Sets and Systems. 58(3): 343-353.
[8] P.K. De and Amita Bhinchar. 2010. Computation of Shortest Path in a fuzzy network. International journal computer applications. 11(2): 0975-8887.
[9] R. Sophia Porchelvi and G. Sudha. 2013. A modified algorithm for solving shortest path problem with Intuitionistic fuzzy arc lengths. International Journal of Scientific and Engineering Research. 4(10): 22295518.
[10] R. Sophia Porchelvi and G. Sudha. 2014. Intuitionistic Fuzzy Critical path in a network. International Conference on Mathematical Methods and Computations Proceedings.
[11]S. Okada and T. Soper. 2000. A Shortest path problem on a network with fuzzy arc lengths. Fuzzy Sets and Systems. 109(1): 129-140.
[12]T.N. Chuang and J.Y. Kung. 2005. The fuzzy shortest path length and the corresponding shortest path in a network. Computers and Operations Research. 32(Np/6): 1428.
[13] W.J. Wang. 1997. New similarity measures on fuzzy sets and on elements. Fuzzy Sets and Systems. 85(3): 305-309.
[14]Liu S.T and C. Kao. 2004. Network flow problems with fuzzy arc lengths. IEEE Transactions on systems, Man and Cybernetics: Part B. 34: 765-769.
[15]K.Atanassov. 1986. Intuitionistic Fuzzy Sets, Fuzzy sets and System. 20(1): 87-96.
[16] Shu MH, Cheng CH and Chang JR. 2006. Using intuitionistic fuzzy sets for fault-tree analysis on printed circuit board assembly, Microelectron Reliab. 46(12): 2139-2148.
[17]Tzung-Nan Chuang and Jung-Yuan Kung. 2006. A new algorithm for the discrete fuzzy shortest path
problem in a network. Applied Mathematics and Computation. 174, pp. 660-668.
[18] Xinfan Wang. 2008. Fuzzy Number Intuitionistic Fuzzy Arithmetic Aggregation Operators. International Journal of Fuzzy systems. 10(2): 134141.
[19]R. Yager. 1986. Paths of least resistance on possibilistic production systems. Fuzzy Sets and Systems. 19, pp.121-132.
[20]T. Gerstenkorn and J. Man'ko. 1991. Correlation of intuitionistic fuzzy sets. Fuzzy Sets and Systems. 44, pp. 39-43.

